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MPI on Multicore, an OpenMP Alternative?

No matter how you cut it, coding for multicore is really just parallel programming.Doug Eadline explains the differences between OpenMP and MPI, when it's smart to use existing code and when it's time to rewrite an application to scale better on multicore systems.

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No matter how you cut it, coding for multicore is really just parallel programming. Once you've realized that, it's time to look at the options, whether your existing codebase will scale, or if you need to rewrite your code and how.

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As stated in <u>The Multicore Programming Challenge</u>, parallel programming can be difficult. It moves the programmer closer to the hardware and further from their application space or problem. Fortunately, people like rocket scientists have been writing

parallel software for quite some time in the HPC (High Performance Computing) sector.

As any good programmer knows, an existing code base can be valuable to current programming projects. First, the possibility of re-using existing code is a major incentive. Also, learning how someone else attacked a similar problem is very valuable.

In the HPC sector, most parallel programs are written using Message Passing Interface (MPI). While MPI is normally used on large computing systems (clusters) it can be also be used on a multicore processor. The "MPI proposition" may seem counter to conventional wisdom as MPI was designed for distributed memory (i.e. each core/processor has it own private memory), whereas OpenMP was designed for shared memory.

The lazy assumption suggests that OpenMP is a better solution because it was designed for shared memory. However, the possibility of re-using an existing MPI code base is worth considering before you spend a month(s) re-inventing the software wheel. Ultimately, the question is really about efficiency. Namely, *How does the performance of MPI compare to OpenMP on a multicore system?*

The answer to this question is important. If I can re-use MPI codes that work *well enough* on multicore, then there is no need to (re)write my application using OpenMP. If, on the the other hand, OpenMP or threads provide scaling benefits sufficient enough to justify re-writing the code, then investing the time in re-coding might be in order.

Although your application(s) are always the ultimate test of hardware, a comparison of the same program written in MPI and OpenMP would be interesting. Fortunately for us, the people at NASA (the rocket science guys) have an interest in such things as well. The venerable <u>NAS Parallel Suite</u> is now available in MPI, OpenMP, Java, and HPF.

This enhancement means a head to head comparison of MPI and OpenMP is possible. (I'll leave the Java and HPF runs as an exercise for the reader). Before we get to the main event however, some background on how OpenMP and MPI differ may be helpful.

OpenMP and MPI Primer



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Because native Pthread programing can be cumbersome, a higher level of abstraction has been developed called OpenMP. As with all higher level approaches, OpenMP sacrifices flexibility for the ease of writing code. At its core, OpenMP uses threads, but the details are hidden from the programmer.

OpenMP is implemented as compiler directives in program comments. Typically, computationally heavy loops are augmented with OpenMP directives that the compiler uses to automatically "thread the loop". This type of approach has the distinct advantage that it may be possible to leave the original program "untouched" (except for comment-directives) and provide simple recompilation for a sequential (non-threaded) version where the OpenMP directives are ignored. (Read the <u>OpenMP Web site</u> to get the complete picture.)

For those who don't follow software trends, but instead rely on the crack *Linux Magazine* columnists to provide them with all the important advances, GCC 4.2 (and later) has support for OpenMP. This is important for the open source crowd, because OpenMP was only available in commercial compilers before GCC 4.2 was released.

GCC 4.2 has not found its way into all distributions, so you may need to <u>download</u> and build it from source if you want to play along with this article. Of course if you have a commercial compiler, it probably already has OpenMP support.

For gcc and gfortran, OpenMP programs can be compiled by including the -fopenmp option. In order to test this new capability, I found an OpenMP version of the ubiquitous matrix multiplication program. I built two versions of the program, one with OpenMP enabled and one without:

\$ gfortran -fopenmp -o matmult_omp matmult.f \$ gfortran -o matmult matmult.f

Then I ran the sequential version on an Intel Core 2 Duo system (two cores):

\$time ./matmult

real 0m9.079s user 0m8.988s sys 0m0.012s

The OpenMP version was run as well. Note that there is a environment variable called *OMP_NUM_THREADS* that will tell OpenMP binaries how many threads to use. If this is not defined, one thread per CPU (core) is used. Ultimately however, the maximum number of threads may be defined by the program. The OpenMP results for two cores is shown below.)

\$ time ./matmult_omp
real 0m4.967s

user 0m9.783s sys 0m0.018s

The OpenMP version reduced the wall clock time by forty five percent. Astute readers may be wondering, why the *user* time is almost double the *real* time. This effect is due to using two cores, i.e. your total CPU time is a sum of the cores your application is uses. As we will see below, the *user* time can be quite a bit higher than the *real* time for eight cores.

In contrast to OpenMP, MPI uses a software library to send data from one process to another. Each process has its own memory space and thus MPI is basically a message copying methodology. In addition, MPI makes no distinction where a process runs. It can run on the same machine or on another machine. If one were to *time* an 8-way OpenMP and MPI program, the following would result (OpenMP is run first.):

\$time bin/cg.B real 1ml1.735s user 9m23.287s sys 0m2.012s \$time mpirun -np 8 bin/cg.B.8 real 1m16.138s user 0m0.000s sys 0m0.004s

In the first case, OpenMP shows a *real* time of about one minute with *user* time of almost 9 and a half minutes indicating a good speed up. In the second case, the MPI run shows a comparable *real* time, but zero *user* time. This result is easily understood in terms of how MPI jobs are run. The *mpirun* command starts each separate MPI process and then waits until they are finished, thus no *user* time. OpenMP jobs, however, share a process space which makes them tractable to the OS.

The Process View

While we are talking about OpenMP and MPI, there's one big difference between these programming methods in terms of the OS process space. OpenMP programs run as a single

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The HP BladeSystem C3000 Same capability. Less space. No compromise. process and the parallelism is expressed as threads. This behavior can be viewed quite clearly when using an eight core server (two quad-core processors). For instance, examining a running OpenMP program using *top* shows only a single process running. (See *Figure One*)

				deadline	@clover	:~/NPI	B 3.2	.1/NP	B3.2	-MPI			_ 0 ×
Eile [<u>E</u> dit <u>V</u> iew <u>T</u> e	ermin	al Ta	bs Help)								
top -	17:05:55	up	7:52	2, 31	users,	108	ad	aver	age	: 2.4	10, 2.34,	1.72	*
Tasks	: 193 tota	1,	11	unning	g, 192	slee	epi	ng,	0	stop	oped, () zombie	
Cpu0	: 98.0%us	, (0.0%s	sy, 0	.0%ni,	2.0)%i	d,	0.09	wa,	0.0%hi,	0.0%si,	0.0%st
Cpu1	: 97.0%us	, (0.7%s	sy, 0	.0%ni,	2.3	3%i	d,	0.0	‰wa,	0.0%hi,	0.0%si,	0.0%st
Cpu2	: 99.0%us	, (0.3%s	sy, 0	.0%ni,	0.7	7%i	d,	0.0	wa,	0.0%hi,	0.0%si,	0.0%st
Cpu3	: 99.0%us	, (0.0%s	sy, 0	.0%ni,	1.0)%i	d,	0.09	wa,	0.0%hi,	0.0%si,	0.0%st
Cpu4	: 98.3%us	, (0.0%s	sy, O	.0%ni,	1.7	7%i	d,	0.09	%wа,	0.0%hi,	0.0%si,	0.0%st
Cpu5	: 98.0%us	, (0.3%s	sy, 0	.0%ni,	1.7	7%i	d,	0.09	‰wа,	0.0%hi,	0.0%si,	0.0%st
Cpu6	: 98.0%us	, i	1.0%s	sy, 0	.0%ni,	1.()%i	d,	0.09	wa,	0.0%hi,	0.0%si,	0.0%st
Cpu7	: 99.0%us	, (D.7%s	sy, O	.0%ni,	0.3	3%i	d,	0.0	%wа,	0.0%hi,	0.0%si,	0.0%st
Mem:	8143108k	to	tal,	5452	124k u	sed,	2	6909	84k	free	e, 2385	52k buffer	s
Swap:	2031608k	to	tal,		0k u	sed,	2	0316	08k	free	e, 43021	52k cached	1
PTD	USER	PR	NT	VTRT	RES	SHR	S	%CPU	1 %M	EM	TTME+	COMMAND	
17701	deadline	18	0	425m	341m	836	S	788	4	.3	2:05.99	cg.B	
3615	root	18	õ	5848	420	320	S	0	0	.0	0:07.87	wulfd	
17695	deadline	15	0	12708	1140	800	R	0	0	.0	0:00.37	top	
1	root	15	0	10304	660	548	S	0	0	.0	0:02.09	init	
2	root	RT	0	0	0	0	S	0	0	.0	0:00.01	migration	0
3	root	34	19	0	0	0	S	0	0	.0	0:00.00	ksoftirad	0
4	root	RT	0	0	0	0	S	0	0	.0	0:00.00	watchdog/()
5	root	RT	0	0	0	0	S	0	0	.0	0:00.00	migration/	1
6	root	34	19	0	0	0	S	0	0	.0	0:00.00	ksoftirqd/	/1
7	root	RT	0	0	0	0	S	0	0	.0	0:00.00	watchdog/1	

Figure One: OpenMP program (cg.B) running on eight cores.

In contrast to the OpenMP, MPI actually starts one process per core using the mpirun -np 8 ... command. This situation is shown in *Figure Two* where an MPI version of the same program is now running. Note the number of processes is now eight. The processor (core) loads are about the same for both, however.

deadline@clover:~/NPB3.2.1/NPB3.2-OMP							×						
<u>E</u> ile [<u>E</u> dit <u>V</u> iew <u>T</u> €	erminal	l Tab	s <u>H</u> elj	2								
top -	17:08:37	up	7:55	, 3	users	, lo	ad	aver	age: 4.	17, 3.10,	2.12		
Tasks	: 201 tota	1,	9 r	unnin	g, 192	2 sle	epi	ng,	0 sto	pped, () zombie		
Cpu0	: 99.0%us	, 1	.0%s	y, 0	.0%ni	. 0.	0%i	d,	0.0%wa,	0.0%hi,	0.0%si,	0.0%st	
Cpu1	: 99.7%us	, 0	.3%s	y, 0	.0%ni	, 0.	0%i	d,	0.0%wa,	0.0%hi,	0.0%si,	0.0%st	
Cpu2	: 99.3%us	, 0	.7%s	y, 0	.0%ni	, 0.	0%i	d,	0.0%wa,	0.0%hi,	0.0%si,	0.0%st	
Cpu3	: 99.3%us	, 0	.7%s	y, 0	.0%ni	0.	0%i	d,	0.0%wa,	0.0%hi,	0.0%si,	0.0%st	
Cpu4	: 99.7%us	, 0	.3%s	y, 0	.0%ni	0.	0%i	d,	0.0%wa,	0.0%hi	0.0%si,	0.0%st	
Cpu5	: 99.0%us	, 1	.0%s	y, 0	.0%ni	0.	0%i	d,	0.0%wa,	0.0%hi,	0.0%si,	0.0%st	
Cpu6	: 99.3%us	, 0	.7%s	y, 0	.0%ni	0.	0%i	d,	0.0%wa,	0.0%hi,	0.0%si,	0.0%st	
Cpu7	:100.0%us	, 0	.0%s	y, 0	.0%ni	0.	0%i	d,	0.0%wa,	0.0%hi,	0.0%si,	0.0%st	
Mem:	8143108k	tot	al,	5540	544k u	used,	2	6025	64k fre	e, 2388	300k buffe:	rs	
Swap:	2031608k	tot	al,		Ok ı	used,	2	0316	08k fre	e, 43039	928k cache	d	
PID	USER	PR	NI	VIRT	RES	SHR	S	%CPU	%MEM	TIME+	COMMAND		
17719	deadline	18	0	112m	51m	2676	R	100	0.7	0:29.23	cg.B.8		
17715	deadline	18	0	112m	59m	3108	R	100	0.8	0:29.25	cg.B.8		
17716	deadline	24	0	112m	58m	2520	R	100	0.7	0:29.26	cg.B.8		
17717	deadline	19	0	112m	51m	2536	R	100	0.6	0:29.25	cg.B.8		
17718	deadline	19	0	112m	51m	2676	R	100	0.7	0:29.25	cg.B.8		
17722	deadline	18	0	112m	59m	2516	R	100	0.7	0:29.23	cg.B.8		
17720	deadline	18	0	112m	51m	2676	R	100	0.6	0:29.17	cg.B.8		
17721	deadline	18	0	112m	59m	2528	R	100	0.7	0:29.20	cg.B.8		
3615	root	18	0	5848	420	320	S	0	0.0	0:08.00	wulfd		
15252	root	15	0	69732	2816	2176	S	0	0.0	0:00.29	sshd		
1	root	15	0	10304	660	548	S	0	0.0	0:02.09	init		•

Figure Two: MPI program (cg.B.8) running on eight cores.

One final and subtle point. In OpenMP communication is through shared memory, which means *threads share access* to a memory location. With MPI programs on SMP systems communication is also through shared memory, but *processes send messages by writing from*

private to shared memory.

Obviously, sharing memory locations seems more efficient than sending copies of memory locations to other processes, but it all depends. In the MPI process model, single processes have exclusive access to all their process memory. For some programs this situation may be more efficient because it is better to copy data (send a message) than to wait for shared memory access. On the other hand, in the OpenMP model, threads can share access to all memory in the process space. In this case, some programs may be much more efficient as the large overhead of copying memory is not needed.

Looking at the Numbers

An eight-core Intel server (two four core Clovertown processors) was used to run the tests. The OpenMP tests used gcc/gfortran version 4.2. The MPI tests used LAM version 7.1.2. The OpenMP and MPI suites have six programs in common and each of these was run five times and averaged (Class B problem sizes were used). The results are given in Mops (million operations per second) in *Table One*. The percent difference is also shown.

Test	OpenMP gcc/gfortran 4.2	MPI LAM 7.1.2	Percent Difference
CG	790.6	739.1	7%
EP	166.5	162.8	2%
FT	3535.9	2090.8	69%
IS	51.1	122.5	139%
LU	5620.5	5168.8	9%
MG	1616.0	2046.2	27%

Table One: Results for the OpenMP/MPI benchmarks. (winning test is in bold)

Tests CG and EP are about the same. Indeed, EP is a good check as both methods should produce a similar result because there is very little communication. OpenMP is the clear winner with FT performance, but MPI does surprisingly better with the latency sensitive IS benchmark. In the fifth test, OpenMP does best with the LU benchmark, while MPI does best with MG. Overall the comparison is a bit of draw.

The results are clear on one point, there is not a definitive winner in this match-up. This result may come as a surprise to those who would assume, OpenMP would easily beat MPI on an multicore machine. (Or any SMP machine for that matter.) Maybe MPI is *good enough* to stand toe-to-toe with OpenMP for many applications.

In only one case (FT), did OpenMP run away from MPI. In other cases, MPI was a clear winner, and taking the time to convert your code to OpenMP would actually result in a performance loss. The story is far from over, more benchmarks are in order using other hardware and commercial compilers.

Other Things to Consider

Getting back to our question, "do I need to re-code my MPI programs for these multicore thingies?," the answer is a resounding maybe not. MPI may just be good enough in many cases. Again, more data, and results for your application are needed for more solid recommendations.

Another important question to ask is how scalable your application is. As more processors are added, parallel execution will always hit a point of diminishing returns. This situation means that creating more threads or processes will not improve performance and it may actually hurt performance. The size of your data set may also come into play. One of the advantages of distributed MPI programs is the ability to distribute large data sets over many processors thereby solving problems that would never fit in an SMP memory space.

If you're considering a writing a new application from scratch, the choice of OpenMP or MPI includes other considerations. OpenMP is designed for shared memory (SMP) machines. As multicore continues to grow the number of processors on an SMP will continue to grow, but OpenMP is not designed to run across multiple machines like MPI.

If you want your application to be portable on clusters and SMP machines, MPI might be the best solution. If, however, you do not envision using more than eight or sixteen cores, then OpenMP is probably one of your best choices if the benchmarks point in that direction. From a conceptual standpoint, those with experience in both paradigms state that using OpenMP and MPI provide a similar learning curve and *nuance level*. There are no shortcuts or free lunches with OpenMP, or MPI for that matter.

Douglas Eadline is the Senior HPC Editor for Linux Magazine.

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